

## Examples of pp-definable subgroups in the homotopy category of a gentle algebra.

This talk is:

- based on arXiv preprint 1911.07691
- the first of two talks, the second based also on 2004.06854

Plan for the talk today:

- motivate via examples from the module category
- discuss model theory of modules
- recall compactly generated triangulated categories] ←
- discuss 'many sorted' model theory for these
- recall 'string complexes' and 'homotopy words'
- give some examples of the groups in the title

In what remains we consider the gentle algebra

$$\Lambda = KQ/I,$$

where

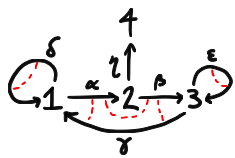
$$\begin{cases} K \text{ is a field} \\ I = \langle \delta\beta, \beta\alpha, \alpha\delta, \delta^2, \varepsilon^2 \rangle \end{cases}$$

{ 1 : pp-definable subgroups of modules.

Consider the word  $w = \beta^{-1}\varepsilon^{-1}\delta^{-1}\delta\delta\varepsilon^{-1}$

2 → 3 → 3 → 1 ← 1 ← 3 → 3

For any  $\Lambda$ -module  $M$  let



$wM = \{m_0 \in e_2 M \mid \exists m_1, \dots, m_6 \in M \text{ such that}$

$$\left( m_i \in \begin{cases} e_i M & (i=3,4) \\ e_3 M & (\text{otherwise}) \end{cases} \right) \text{ and } \left( \begin{matrix} \beta m_0 = m_1, \epsilon m_1 = m_2, \gamma m_2 = m_3 \\ m_3 = \delta m_4, m_4 = \gamma m_5, \epsilon m_5 = m_6 \end{matrix} \right) \}$$

The  $m_i$ 's are described as 'filling in' the schema

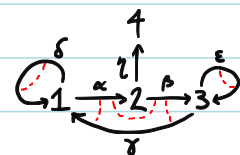
$$m_0 \xrightarrow{\beta} m_1 \xrightarrow{\epsilon} m_2 \xrightarrow{\gamma} m_3 \xleftarrow{\delta} m_4 \xleftarrow{\gamma} m_5 \xrightarrow{\epsilon} m_6$$

Remarks:

- The subspaces  $wM$  arise when one uses 'functorial filtrations' to classify f.d modules [Butler, Ringel, '87]
- if  $\theta: M \rightarrow N$  is a homomorphism of  $\Lambda$ -modules then  $\theta(wM) \subseteq wN$ , hence  $w: \Lambda\text{-Mod} \rightarrow \text{Vect}_k$  is a functor

We can realise elements of  $wM$  by means of  $\Lambda$ -module homomorphisms  $\theta: S \rightarrow M$  where  $S$  is a string module.

To see this, let  $S = \bigoplus_{n=0}^{11} K b_n$  and define the action of  $\Lambda$  using the diagram



$$b_0 \xleftarrow{\delta} b_1 \xrightarrow{\beta} b_2 \xrightarrow{\epsilon} b_3 \xrightarrow{\gamma} b_4 \xleftarrow{\delta} b_5 \xleftarrow{\gamma} b_6 \xrightarrow{\epsilon} b_7 \xrightarrow{\gamma} b_8 \xrightarrow{\delta} b_9 \xrightarrow{\alpha} b_{10} \xrightarrow{\beta} b_{11}$$

$$m_0 \xrightarrow{\beta} m_1 \xrightarrow{\epsilon} m_2 \xrightarrow{\gamma} m_3 \xleftarrow{\delta} m_4 \xleftarrow{\gamma} m_5 \xrightarrow{\epsilon} m_6$$

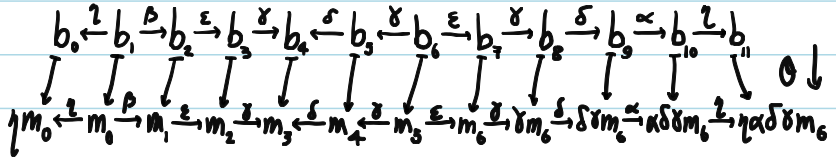
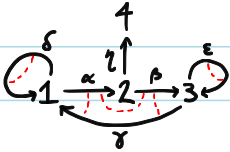
So  $b_0 = \gamma b_1$ ,  $\beta b_1 = b_2$ , and so on. When  $M$  is an arrow such that  $b_i$  is not the tail of  $M$ ,  $M b_i = 0$ . So, for example,  $\epsilon b_7 = 0$ . There are also relations such as  $e_+ b_0 = b_0$ ,  $e_2 b_1 = b_1$ , and so on. Note this gives  $\epsilon^2 b_6 = \epsilon b_7 = 0, \dots$   
We now claim

$$wM = \{ \theta(b_i) : \theta \in \text{Hom}_{\Lambda\text{-Mod}}(S, M) \}$$

$$w = \beta^{-1} \epsilon^{-1} \gamma^{-1} \delta \gamma \epsilon^{-1}$$

For  $\subseteq$ , given  $m_0 \in wM$  choose  $m_1, \dots, m_6$  as before.  
 Then define  $\theta$  by

$$m_0 \xrightarrow{\beta} m_1 \xrightarrow{\varepsilon} m_2 \xrightarrow{\gamma} m_3 \xrightarrow{\delta} m_4 \xrightarrow{\gamma} m_5 \xrightarrow{\varepsilon} m_6$$



The inclusion  $\supseteq$  is easier.

We now interpret  $wM$  using model theory (logic).  
 We claim there exists  $A \in \text{Mat}_{15 \times 7}(\Omega)$  such that

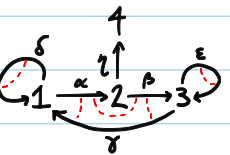
$$wM = \left\{ m_0 \in M : \exists m_1, \dots, m_6 \in M \text{ with } A \begin{pmatrix} m_0 \\ \vdots \\ m_6 \end{pmatrix} = \begin{pmatrix} 0 \\ \vdots \\ 0 \end{pmatrix} \right\}$$

Proof:

$$w = \beta^{-1} \varepsilon^{-1} \delta^{-1} \gamma \delta \varepsilon^{-1}$$

Choose  $A =$

$1 - e_2$	0	0	0	0	0	0	$e_2 m_0 = m_0$
0	$1 - e_3$	0	0	0	0	0	$e_3 m_1 = m_1$
0	0	$1 - e_3$	0	0	0	0	$\vdots$
0	0	0	$1 - e_1$	0	0	0	$\vdots$
0	0	0	0	$1 - e_1$	0	0	$\vdots$
0	0	0	0	0	$1 - e_3$	0	$\vdots$
0	0	0	0	0	0	$1 - e_3$	$e_3 m_6 = m_6$
$\beta$	$-e_3$	0	0	0	0	0	$\beta m_0 = m_1$
0	$\varepsilon$	$-e_3$	0	0	0	0	$\varepsilon m_1 = m_2$
0	0	$\gamma$	$-e_1$	0	0	0	$\vdots$
0	0	0	$-e_1$	$\delta$	0	0	$\vdots$
0	0	0	0	$-e_1$	$\gamma$	0	$\vdots$
0	0	0	0	0	$\varepsilon$	$-e_3$	$\varepsilon m_5 = m_6$



$$m_0 \xrightarrow{\beta} m_1 \xrightarrow{\varepsilon} m_2 \xrightarrow{\gamma} m_3 \xrightarrow{\delta} m_4 \xrightarrow{\gamma} m_5 \xrightarrow{\varepsilon} m_6$$

Note that we could replace  $A$  with a smaller matrix, but for this talk we only care that some  $A$  works.

The last remark above says that  $\text{WM}$  is the set of solutions to a positive primitive (=pp) formula.

In general: a pp-formula (in one free variable  $x_0$ ) has the form

$$\varphi(x_0) = \left( \exists x_1, \dots, x_{c-1} : \bigwedge_{i=1}^r \sum_{j=0}^{c-1} a_{ij} x_j = 0 \right)$$

where  $(a_{ij})_{i,j} \in \text{Mat}_{r \times c}(\Lambda)$  ( $r, c \geq 1$ ).

*ignore when  $c=1$*

This is a formula in the language of  $\Lambda$ -modules.

## § 2: model theory of modules

The language of modules over  $\Lambda$  is defined and denoted

$$\mathcal{L}_\Lambda = \langle \text{pred}_\Lambda, \text{func}_\Lambda, \text{ar}_\Lambda \rangle$$

where

$\text{pred}_\Lambda = \{0\}$ ,  $\text{func}_\Lambda = \{+\} \cup \{a- : a \in \Lambda\}$   
are the predicate and function symbols

The arity function  $\text{ar}_\Lambda : \text{pred}_\Lambda \cup \text{func}_\Lambda \rightarrow \mathbb{N}$  is given by  
 $0 \mapsto 0, + \mapsto 2, a- \mapsto 1$

By introducing a countable set of variables  $x_i$  ( $i \in \mathbb{N}$ ), the terms inside  $\mathcal{L}_\Lambda$  are defined by saying:

- each  $x_i$  is a term
- if  $t$  and  $t'$  are terms then  $t+t'$  is a term
- if  $t$  is a term and  $a \in \Lambda$  then  $at$  is a term

The atomic formulas are those of the form

$$t = t' \text{ or } t = 0$$

where  $t$  and  $t'$  are terms. The formulas in  $\mathcal{L}_\Lambda$  are built from: the variables  $x_i$ , atomic formulas, binary connectives  $\wedge$ ,  $\vee$ , and  $\Rightarrow$ , negation  $\neg$ , and the quantifiers  $\forall$  and  $\exists$ .

A positive primitive formula  $\varphi(x_0, \dots, x_{n-1})$  in  $n \geq 1$  free variables has the form

$$\varphi(x_0, \dots, x_{n-1}) = (\exists y_1, \dots, y_m \bigwedge_{i=1}^r \psi_i(x_0, \dots, x_{n-1}, y_1, \dots, y_m))$$

where each  $\psi_i(x_0, \dots, x_{n-1}, y_1, \dots, y_m)$  is atomic.

The  $\mathcal{L}_\Lambda$ -structures are given by a set  $M$  together with an interpretation  $O_m \in M$  of the predicate symbol, and of the functions  $+_m: M^2 \rightarrow M$  and  $a_m^-: M \rightarrow M$  ( $a \in \Lambda$ ).

The theory  $\text{Th}_\Lambda$  of  $\Lambda$ -modules is given by the usual module axioms, for example distributivity is given by the formula

$$\forall x_0, x_1, a(x_0 + x_1) = a(x_0) + a(x_1) \quad (a \in \Lambda)$$

A model for  $\mathcal{T}h_\Lambda$  is an  $\mathcal{L}_\Lambda$ -structure satisfying the module axioms: hence is the same thing as a  $\Lambda$ -module. There is also a notion of a homomorphism of  $\mathcal{L}_\Lambda$ -structures, which corresponds to a homomorphism of  $\Lambda$ -modules.

Positive primitive formulas are important in the model theory of modules due to the following.

Theorem (Baur, '76) Any formula in  $\mathcal{L}_\Lambda$  is equivalent (relative to  $\mathcal{T}h_\Lambda$ ) to a boolean combination ( $\wedge, \vee, \neg$ ) of  $\forall \exists$  sentences and positive primitive formulas.

{3: compactly generated  $\Lambda$ -ated categories.

Let  $\mathcal{T}$  be a triangulated category with suspension functor  $\Sigma: \mathcal{T} \rightarrow \mathcal{T}$ . The only example we are interested in today is

$$\mathcal{T} = \mathcal{K}(\Lambda\text{-Proj})$$

$\swarrow$  complexes of  $\Lambda$ -projective modules.  
 $\nwarrow$  homotopy category.

\* possibly inf. dim.

We assume also  $\mathcal{T}$  has small coproducts.

So if  $I$  is any set and  $X, Y_i$  ( $i \in I$ ) are objects in  $\mathcal{T}$ , the universal properties give a morphism

$$\coprod_i \text{Hom}_{\mathcal{T}}(X, Y_i) \rightarrow \text{Hom}_{\mathcal{T}}(X, \coprod_i Y_i)$$

We call the object  $X$  compact provided, for any such collection  $(Y_i : i \in I)$ , the map above is bijective.

We say  $\mathcal{T}$  is compactly generated by a set  $\mathcal{G}$  of objects in  $\mathcal{T}$  provided:

- each  $X \in \mathcal{G}$  is a compact object in  $\mathcal{T}$ ,
- if  $X \in \mathcal{G}$  then  $\Sigma X \in \mathcal{G}$ , and
- if  $M$  is an object in  $\mathcal{T}$  with  $\text{Hom}_{\mathcal{T}}(X, M) = 0$  for all  $X \in \mathcal{G}$ , then we must have  $M = 0$ .

← generalised by Neeman '08

Theorem (Jørgensen '05) Since  $\Lambda$  is sufficiently nice, the category  $\mathcal{K}(\Lambda\text{-Proj})$  is compactly generated.

Later we will combine details from this theorem with results from [BT, Thesis], to describe a set of isoclasses of indecomposable compact objects in  $\mathcal{K}(\Lambda\text{-Proj})$ .

Assumption: in what follows we assume that  $\mathcal{T}$  is compactly generated, (that  $\mathcal{T}$  has small coproducts) and that the full subcategory  $\mathcal{T}^c$  of  $\mathcal{T}$  consisting of compact objects is skeletally small.

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§ 4: The canonical many-sorted language.

Remarks:

- in the second talk we'll look at notions of purity, adapted from  $\Lambda\text{-Mod}$  to  $\mathcal{T}$  by Krause
- for today we replace  $\mathcal{L}_{\Lambda}$  by the following

Definition (Garkusha, Prest, '05) Let  $\mathcal{T}$  be a triangulated category satisfying the above Assumption. Fix a chosen set  $\mathcal{S}$  of isoclasses inside  $\mathcal{T}^c$

The canonical  $\mathcal{S}$ -sorted language for  $\mathcal{T}$  is defined and denoted

$$\mathcal{L}_{\mathcal{T}}^{\mathcal{S}} = \langle \text{pred}_{\mathcal{T}}^{\mathcal{S}}, \text{func}_{\mathcal{T}}^{\mathcal{S}}, \text{ar}_{\mathcal{T}}^{\mathcal{S}}, \text{sort}_{\mathcal{T}}^{\mathcal{S}} \rangle$$

where

$$- \text{pred}_{\mathcal{T}}^{\mathcal{S}} = \{0_G : G \in \mathcal{S}\},$$

$$- \text{func}_{\mathcal{T}}^{\mathcal{S}} = \{+_G : G \in \mathcal{S}\} \cup \{o_a : a \in \text{Hom}_{\mathcal{T}}(G, H), G, H \in \mathcal{S}\}$$

$$- \text{ar}_{\mathcal{T}}^{\mathcal{S}}(0_G) = 0, \text{ar}_{\mathcal{T}}^{\mathcal{S}}(+_G) = 2, \text{ar}_{\mathcal{T}}^{\mathcal{S}}(o_a) = 1$$

-  $\text{sort}_{\mathcal{T}}^{\mathcal{S}} : \text{pred}_{\mathcal{T}}^{\mathcal{S}} \cup \text{func}_{\mathcal{T}}^{\mathcal{S}} \rightarrow \bigsqcup_{n \in \mathbb{N}} \mathcal{S}^n$  is a function given by

$$0_G \mapsto G \in \mathcal{S}^1, +_G \mapsto (G, G, G) \in \mathcal{S}^3, o_a \mapsto (H, G) \in \mathcal{S}^2$$

As we did with the language  $\mathcal{L}_{\Lambda}$  of  $\Lambda$ -modules, one can define sorted terms, atomic formulas, formulas and pp-formulas in  $\mathcal{L}_{\mathcal{T}}^{\mathcal{S}}$ . For example one begins by fixing a countable set  $\mathcal{V}_G$  of variables for each sort  $G \in \mathcal{S}$ , each term has a sort, ...

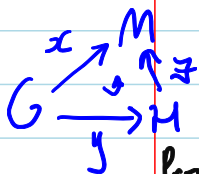
To understand the above... Given any object  $M$  in  $\mathcal{T}$ , we define an  $\mathcal{L}_{\mathcal{T}}^{\mathcal{S}}$ -structure  $\underline{M}$  as follows.

Let  $\underline{M} = \{ \text{Hom}_{\mathcal{T}}(G, M) : G \in \mathcal{S} \}$ , where

- the symbol  $0_G$  is interpreted as  $0 \in \text{Hom}_{\mathcal{T}}(G, M)$
- the symbol  $+_G$  ——— " ——— the addition here
- each symbol  $o_a$  ——— " ———  $\text{Hom}_{\mathcal{T}}(H, M) \rightarrow \text{Hom}_{\mathcal{T}}(G, M)$   
 $f \mapsto fa$



Of great importance later will be the pp-formulas given by divisibility; that is those of the form



$$\varphi(x_G) = (\exists y_H: x_G = y_H \circ a)$$

$G \rightarrow M \qquad H \rightarrow M \qquad G \rightarrow M$

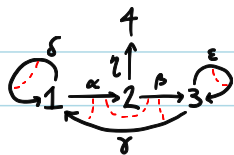
Proposition (Garkusha, Prest) Every pp-formula in the canonical language  $L_T^s$  is equivalent to a pp-formula of the form above (divisibility condition).

Definition (Garkusha, Prest) For any object  $M$  in  $T$  and any object  $G$  in  $T^c$ , a pp-definable subgroup of  $M$  of sort  $G$  is a subgroup of  $\text{Hom}_T(G, M)$  of the form

$$M a = \{ f \circ a \mid f \in \text{Hom}_T(H, M) \}$$

where  $a: G \rightarrow M$  is some morphism in  $T^c$ .

{ 5: pp-definable subgroups in  $K(\Lambda\text{-Proj})$



Back to the gentle algebra  $\Lambda = KQ/I$ . We want to parallel the story at the start of the talk, which was about  $\Lambda\text{-Mod}$ .

Write  $K_{\min}(\Lambda\text{-Proj})$  for the full subcategory of  $K(\Lambda\text{-Proj})$  consisting of complexes  $M^\bullet$  such that  $\text{im}(d_n^\bullet) \subseteq \text{rad}(M^{n+1})$  for all  $n \in \mathbb{Z}$ .

Fact: since  $\Lambda$  is a perfect ring the subcategory  $K_{\min}(\Lambda\text{-Proj})$  is dense inside  $K(\Lambda\text{-Proj})$ .

Fix a complex  $M^\bullet$  in  $K_{\min}(\Lambda\text{-Proj})$ . For all  $m \in M = \bigoplus_{n \in \mathbb{Z}} M^n$  we have, for each vertex  $v$  in  $Q$ ,

$$d_m(e_v m) \in e_v \text{rad}(M) = e_v \text{rad}(\Lambda) M,$$

definition of  $K_{\min}$        $M$  is projective

and so

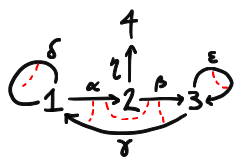
$$d_m(e_v m) \in \begin{cases} \delta M \oplus \gamma M & (v=1) \\ \alpha M & (v=2) \\ \beta M \oplus \varepsilon M & (v=3) \\ \eta M & (v=4) \end{cases}$$

Hence if  $\lambda$  is an arrow in  $Q$  with head  $v$  there is a  $K$ -linear map  $d_{\lambda, m}: e_v M \rightarrow e_v M$  given by

$$\begin{array}{ccc} e_v M & \xrightarrow{d_{\lambda, m}} & e_v M \\ d_m|_{e_v m} \downarrow & \circlearrowleft & \uparrow \cong \\ e_v \text{rad}(M) & \longrightarrow & \lambda M \end{array}$$

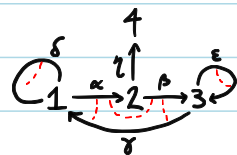
For example, if  $d_m(e_2 m) = \delta l + \gamma n$  then we have  $d_{\delta, m}(m) = \delta l$  and  $d_{\gamma, m}(m) = \gamma n$ .

We can construct homotopy words using the letters of the form  $d_\lambda$ ,  $d_\lambda^{-1}$ ,  $p$  and  $p^{-1}$  where  $p$  is a non-zero path in  $\Lambda$  and  $\lambda$  is an arrow.



For example,  $C = d_s^{-1} \delta \gamma \varepsilon^{-1} \gamma^{-1} \delta^{-1} \alpha^{-1} d_\alpha \beta^{-1} d_\beta d_\varepsilon^{-1} \varepsilon$  is a homotopy word.   
legit because  $\beta\alpha = 0$  in  $\Lambda$   
homotopy letters

For any object  $M$  of  $K_{\min}(\Lambda\text{-Proj})$  let  $(M = \bigoplus_{n \in \mathbb{Z}} M_n)$



$$CM = \{m_0 \in e_1 M \mid \exists m_1 \in e_2 M, m_2 \in e_2 M, m_3 \in e_3 M, m_4 \in e_3 M \text{ such that } \begin{cases} d_{s,M}(m_0) = \delta \gamma m_1, & \alpha \delta \varepsilon m_1 = d_{\alpha,n}(m_2), \\ \beta m_2 = d_{\beta,M}(m_3), & \text{and } d_{\varepsilon,M}(m_3) = \varepsilon m_4 \end{cases} \}$$

Here the  $m_i$ 's fill in the schema

$$m_0 \xrightarrow{d_s} \xleftarrow{\delta \gamma} m_1 \xrightarrow{\alpha \delta \varepsilon} \xleftarrow{d_\alpha} m_2 \xrightarrow{\beta} \xleftarrow{d_\beta} m_3 \xrightarrow{d_\varepsilon} \xleftarrow{\varepsilon} m_4$$

Aim for the remainder today: show that there are compact objects  $G, H$  and a morphism  $a: G \rightarrow H$  in  $K_{\min}(\Lambda\text{-Proj})$  such that the vector space  $CM$  is given by the pp-definable subgroup

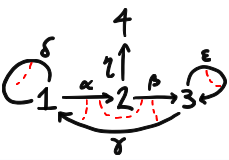
$$Ma = \{ f \circ a \mid f \in \text{Hom}_{K_{\min}(\Lambda\text{-Proj})}(H, M) \}$$

(of  $M$  of sort  $G$ ).

We will use the language of string complexes

The algebras considered here are derived-discrete, but their argument generalises easily.

Lemma (Arnesen, Laking, Paukszetello, Prest '17) Using the Theorem above, due to Jørgensen, the indecomposable compact objects in  $K_{\min}(\Lambda\text{-Proj})$  are given by (shifts of) string and band complexes in  $K_{\min}^{+,b}(\Lambda\text{-proj})$    
 $G = 0, n < 0$   
 $H^i(G) = 0, n > 0$

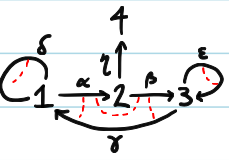
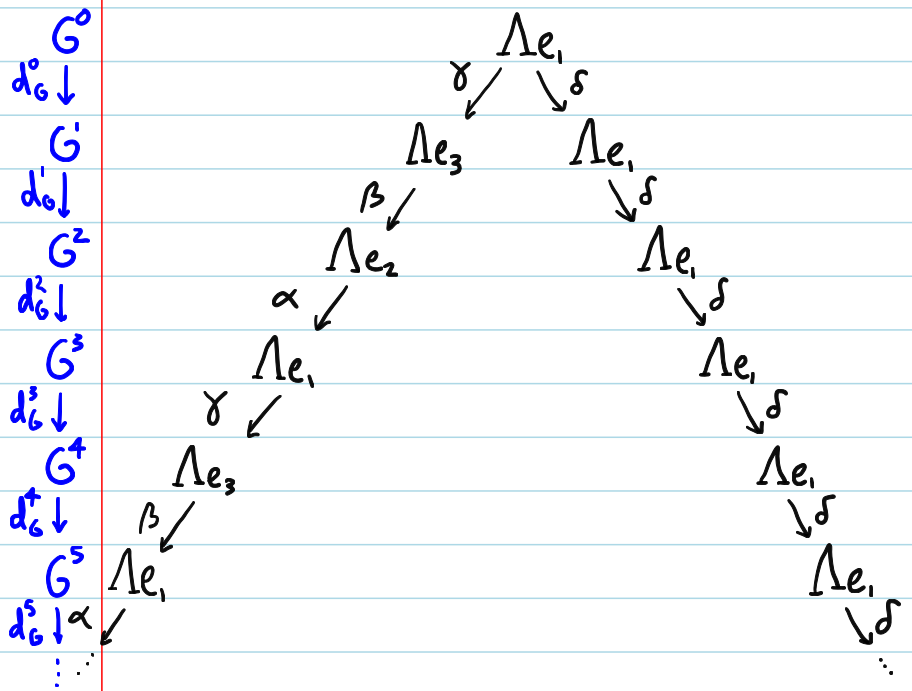


Define the (string) complex  $G^\bullet$  in  $K_{\min}^{+,b}(\Lambda\text{-proj})$  by

$$\begin{array}{ccccccc} & & \begin{pmatrix} \delta \\ \delta \end{pmatrix} & \begin{pmatrix} \beta & 0 \\ 0 & \delta \end{pmatrix} & \begin{pmatrix} \alpha & 0 \\ 0 & \delta \end{pmatrix} & \begin{pmatrix} \gamma & 0 \\ 0 & \delta \end{pmatrix} & \dots \\ \dots \rightarrow & 0 \rightarrow & \Lambda e_1 \rightarrow & \Lambda e_3 \oplus \Lambda e_1 \rightarrow & \Lambda e_2 \oplus \Lambda e_1 \rightarrow & \Lambda e_1 \oplus \Lambda e_1 \rightarrow & \Lambda e_3 \oplus \Lambda e_1 \rightarrow \dots \\ & & \parallel & \parallel & \parallel & \parallel & \parallel \\ \dots \rightarrow & G^{-1} \rightarrow & G^0 \rightarrow & G^1 \rightarrow & G^2 \rightarrow & G^3 \rightarrow & G^4 \rightarrow \dots \end{array}$$

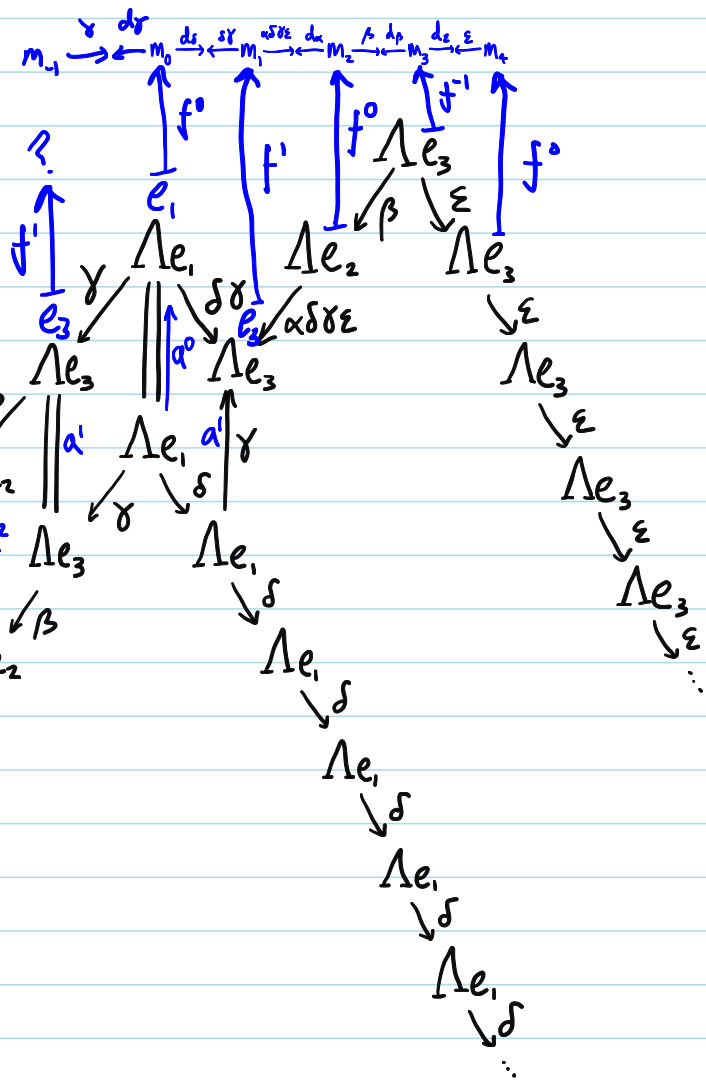
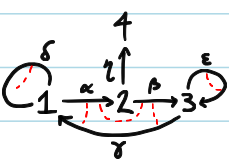
Here we abuse notation, writing  $\gamma$  for the  $\Lambda$ -module homomorphism  $-\gamma: \Lambda e_1 \rightarrow \Lambda e_3, z \mapsto z\gamma$ .

The name string complex is motivated by the redrawing of  $G^\bullet$  using the diagram below.



So far,  $G^\bullet$  seems unrelated to the homotopy word  $C$ . The only connection is the 'head' of  $C$  is 1, and  $G^\bullet$  is the  $\text{Hom}(-, \Lambda)$ -dual of a min. res. of  $e_1 \Lambda / \text{rad}(e_1 \Lambda)$ .

Using the depiction of string complexes above, we finally define the morphism  $a: G \rightarrow H$  in  $K_{\min}^{+,b}(\Lambda\text{-proj})$  such that CM is given by



$$M_a = \{f \circ a \mid f \in \text{Hom}_T(H, M)\}.$$

