- **1.** (a) Let R be a boolean ring and  $u, v, x, y, z \in R$ .
  - (i)  $\wedge = \cdot$  and + are associative by definition of a ring and together with the law of distributivity and the commutativity of + we compute for  $\vee$ :

 $(x \lor y) \lor z = x + y + z + xy + xz + yz + xyz = x \lor (y \lor z)$ 

(ii)  $\wedge = \cdot$  is idempotent by definition of a boolean ring and using 2 = 0 we get:

$$x \lor x = 2x + x^2 = 0 + x = x$$

- (iii) + is commutative by definition of a ring and  $\wedge = \cdot$  by Lemma 1.1.3. This implies the commutativity of  $\lor$  and  $\leftrightarrow$ , which in turn yields that of  $\uparrow$  and  $\downarrow$ .
- (iv) Clearly,  $x \land 0 = 0$ ,  $x \land 1 = x$ ,  $x \lor 0 = x$  and  $x \lor 1 = 2x + 1 = 1$ , since 2 = 0.
- (v) Using the idempotence of  $\wedge = \cdot$  and 2 = 0:

$$x \wedge (x \vee y) = x(x + y + xy) = x + 2xy = x$$
$$x \vee (x \wedge y) = x + xy + x^2y = x + 2xy = x$$

(vi) Using 2 = 0:

$$\neg (x \land y) = xy + 1 + 2(x + y + 1) = \neg x \lor \neg y$$
  
$$\neg (x \lor y) = x + y + xy + 1 = \neg x \land \neg y$$

(vii) Using  $\neg \neg y = y$  (because of 2 = 0) and the commutativity of  $\lor$ :

 $x \to y = \neg x \lor y = \neg \neg y \lor \neg x = \neg y \to \neg x$ 

(viii) Using the idempotence of  $\wedge = \cdot$  and 2 = 0 we compute

$$x \leftrightarrow y = (\neg x \lor y)(\neg y \lor x) = (xy + x + 1)(xy + y + 1) = x + y + 1,$$
  
such that  $x \leftrightarrow y = 1 \Leftrightarrow x + y = 0 \Leftrightarrow x = y.$ 

(ix) Using (vi) and the associativity of  $\lor$ :

$$(x \land y) \to z = \neg (x \land y) \lor z = \neg x \lor \neg y \lor z = x \to (y \to z)$$

(x) Using the idempotence of  $\wedge = \cdot$  and 2 = 0:

$$(x \land \neg y) \lor (\neg x \land y) = (xy + x) \lor (xy + y) = x + y + 6xy = x + y (x \lor y) \land \neg (x \land y) = (x + y + xy)(xy + 1) = x + y + 4xy = x + y$$

(xi) Using  $x \to y = xy + x + 1$ ,  $x \land y = xy$ ,  $x \lor y = xy + x + y$ ,  $x \land \neg y = xy + x$ and 2 = 0 it is easy to see:

$$x \leq y \Leftrightarrow x \rightarrow y = 1 \Leftrightarrow x \land y = x \Leftrightarrow x \lor y = y \Leftrightarrow x \land \neg y = 0$$

Clearly,  $xy = x \Rightarrow y \mid x$ . Conversely, let's assume now x = uy for some u. From (xuy)xy = xuy and (xy)x = xy we get  $x = xuy \le xy \le x$ . Hence, we have  $x \leftrightarrow xy = (x \to xy)(xy \to x) = 1$  such that x = xy by (viii), so  $x \le y$ .

- (xii) Let  $u \leq x$  and  $v \leq y$ . By (xi) this means  $u \wedge x = u$  and  $v \wedge y = v$  such that  $(u \wedge v) \wedge (x \wedge y) = u \wedge v$  by associativity of  $\wedge$ , i.e.  $u \wedge v \leq x \wedge y$ . Similarly,  $u \vee x = x$  and  $v \vee y = y$  such that  $(u \vee v) \vee (x \vee y) = x \vee y$ , i.e.  $u \vee v \leq x \vee y$ .
- (xiii) If  $x \wedge y = 0$  and  $x \vee y = 1$ , then  $y = x + x \wedge y + x \vee y = x + 1 = \neg x$ .

- (b) Let R be a boolean ring. Then we have:
  - (I) (xi) and (iv) yield  $0 \le x \le 1$ . We have  $x \le x$  because of  $x \land x = x$  in view of (ii) and (xi). If  $x \le y$  and  $y \le x$ , then  $x \leftrightarrow y = 1$ , so x = y by (viii). If  $x \le y \le z$ , then  $x = x \lor 0 \le y \lor z = z$  by (iv), (xii), and (xi).
  - (II) By (v) and (xi)  $x \wedge y$  is a lower bound and  $x \vee y$  an upper bound of  $\{x, y\}$ . If  $z \leq x$  and  $z \leq y$ , then  $z = z \wedge z \leq x \wedge y$  by (xii) and (ii). If  $x \leq z$  and  $y \leq z$ , then  $x \vee y \leq z \vee z = z$  by (xii) and (ii).
  - (III) Using the idempotence of  $\wedge = \cdot$  and 2 = 0:

$$x \wedge (y \lor z) = xy + xz + xyz = (x \land y) \lor (x \land z) x \lor (y \land z) = x + yz + xyz = (x \lor y) \land (x \lor z)$$

(IV) Using the idempotence of  $\wedge = \cdot$  and 2 = 0:

$$x \wedge \neg x = x(x+1) = 2x = 0 x \vee \neg x = x + (x+1) + x(x+1) = 4x + 1 = 1$$

(c) Let  $(R, \leq 0, 1, \neg, \land, \lor)$  be a boolean algebra and define  $\cdot = \land$  and + as follows:

$$x + y = (x \land \neg y) \lor (\neg x \land y)$$

It is clear that  $\cdot = \wedge$  and  $\vee$  are commutative and idempotent. They are associative, since for  $\diamond = \wedge$  resp.  $\diamond = \vee$  both  $(x \diamond y) \diamond z$  and  $x \diamond (y \diamond z)$  are easily checked to be the infimum resp. supremum of  $\{x, y, z\}$ . We will use the following observations:

## Lemma.

(1) If  $x \wedge y = 0$  and  $x \vee y = 1$ , then  $y = \neg x$ . (2)  $\neg (x + y) = (\neg x \wedge \neg y) \vee (x \wedge y)$ . (3)  $\neg (x \wedge y) = \neg x \vee \neg y$ .

*Proof.* (1) We calculate  $y = y \land 1 = y \land (x \lor \neg x) = (x \land y) \lor (y \land \neg x) = y \land \neg x \le \neg x$ . Interchanging the roles of y and  $\neg x$ , we get  $\neg x \le y$ . Hence,  $y = \neg x$ .

For (2) choose  $z = (\neg x \land \neg y) \lor (x \land y)$  and w = x + y. For (3) choose  $z = \neg x \lor \neg y$ and  $w = x \land y$ . In both cases  $z \land w = 0$  and  $z \lor w = 1$  such that we can use (1).  $\Box$ 

Now it is easy to see that  $(R, +, \cdot)$  is a ring:

- The associativity of  $\cdot$  and  $1 \cdot x = x = x \cdot 1$  show that  $(R, \cdot)$  is a monoid.
- The commutativity of  $\land$  and  $\lor$  yields the commutativity of +. Since  $\neg 0 = 1$  by (1), we have  $x + 0 = (x \land 1) \lor (\neg x \land 0) = x \lor 0 = x$ . By a straightforward calculation using (2) both (x + y) + z and x + (y + z) are equal to

$$(x \wedge \neg y \wedge \neg z) \lor (\neg x \wedge y \wedge \neg z) \lor (\neg x \wedge \neg y \wedge z) \lor (x \wedge y \wedge z),$$

proving the associativity of +. Finally,  $x + x = 0 \lor 0 = 0$ , i.e. each element of R is its own inverse w.r.t. +. This shows that (R, +) is an abelian group.

• Using (3) for the first identity we get the distributive law

$$\begin{array}{l} xy + xz = ((x \land y) \land (\neg x \lor \neg z)) \lor ((\neg x \lor \neg y) \land (x \land z)) \\ = & (x \land y \land \neg z) \lor (x \land \neg y \land z) \\ \end{array} = x(y+z) + x($$

**2.** Clearly,  $0, 1 \in R'$ . Let  $x, y, z \in R'$ . The calculation

 $(x-y)^4 = x - 4xy + 6xy - 4xy + y = x - 2xy + y = (x-y)^2$ 

shows that  $x + y = (x - y)^2 = x + y - 2xy$  yields an operation +' on R'. It has 0 as neutral element, satisfies x + x = 0, is commutative and also associative because of

$$(x + 'y) + 'z = x + y + z - 2(xy + xz + yz) + 4xyz = x + '(y + 'z).$$

Thus (R', +') is an abelian group. Given that  $\cdot$  is commutative, we see that  $x \cdot y = x \cdot y$  induces an operation  $\cdot'$  on R'. It only remains to check the distributive law:

$$x \cdot '(y + 'z) = x(y - z)^{2} = (xy - xz)^{2} = x \cdot 'y + 'x \cdot 'z$$

**3.** We have  $D_{xy} = D_x \cdot D_y$ , since  $xy \notin \mathfrak{p} \Leftrightarrow (x \notin \mathfrak{p} \text{ and } y \notin \mathfrak{p})$  for prime ideals  $\mathfrak{p}$ .

For  $D_{x+y} = D_x + D_y$  we have to verify  $x + y \notin \mathfrak{p} \Leftrightarrow (x \notin \mathfrak{p}, y \in \mathfrak{p} \text{ or } y \notin \mathfrak{p}, x \in \mathfrak{p})$ for all prime ideals  $\mathfrak{p}$ . The implication  $\Leftarrow$  is true, since  $\mathfrak{p}$  is an ideal. To check  $\Rightarrow$ assume now  $x + y \notin \mathfrak{p}$ . If  $x \in \mathfrak{p}$ , then necessarily  $y \notin \mathfrak{p}$  because  $\mathfrak{p}$  is an ideal, so we are done. If we had  $x \notin \mathfrak{p}$  and  $y \notin \mathfrak{p}$ , we would have  $\neg x, \neg y \in \mathfrak{p}$  by Lemma 1.1.4, since  $\mathfrak{p}$  is a prime ideal, and thus the contradiction  $x + y = \neg x + \neg y \in \mathfrak{p}$ .

Clearly,  $D_1 = \text{Spec}(R) = 1$  because prime ideals are proper.

This shows that D is a ring homomorphism. It remains to verify its injectivity. So let  $x \in R$  with  $x \neq 0$ . Then  $\neg x$  is a zero divisor (because of  $\neg x \land x = 0$ ) and thus contained in some maximal ideal  $\mathfrak{p}$  of R. By Lemma 1.1.4 we have  $x \notin \mathfrak{p}$ , so  $D_x \neq 0$ .

The final statement readily follows from  $\mathcal{P}(X) \cong {}^X\mathbb{F}_2$  (see Remark 1.1.9).

**4.** Let us write  $\circ$ , -,  $\setminus$  for the functions that associate with subsets  $U \subseteq X$  their interior  $\circ U = U^{\circ}$ , their closure  $-U = \overline{U}$ , and their complement  $\setminus U = X \setminus U$ .

From general topology we have the following properties for all  $U, V \subseteq X$ :

- (i)  $\circ$  and preserve  $\subseteq$  and we have  $\circ \subseteq$  pointwise.
- (ii)  $\circ \circ = \circ$  and --=- and  $\setminus -= \circ \setminus$  and  $\setminus \circ = \setminus$ .
- (iii)  $\circ(U \cap V) = \circ U \cap \circ V$  and  $-(U \cup V) = -U \cup -V$ .
- (iv)  $\circ U \cap -V \subseteq -(U \cap V)$ .

For the possibly non-standard fact (iv) just observe that for open M with  $M \subseteq U$ and closed N with  $U \cap V \subseteq N$  the set  $\backslash (M \cap \backslash N)$  is closed and contains V.

## **Lemma.** For all $U, V \subseteq X$ :

(1) 
$$U$$
 regular  $\Leftrightarrow U = \neg \neg U$ .  
(2)  $\circ \subseteq \neg \neg$ .  
(3)  $\neg \neg \circ = \neg \circ$ .  
(4)  $\neg \neg = \neg \neg$ .  
(5)  $\circ U \cap \neg \neg V \subseteq \neg \neg (U \cap V)$ .  
(6)  $\neg \neg (U \cap V) \subseteq \neg \neg U \cap \neg \neg V$  with equality if either U or V is open.

*Proof.* (1) Since  $\neg = \backslash - = \circ \backslash$  we have  $\neg \neg = \circ \backslash \backslash - = \circ -$ .

(2) Applying  $\-\$  on the left to  $\circ \subseteq -$  gives  $\circ = \circ \circ = \circ \\) \circ = \-\) \circ \subseteq \-\-\- = \neg \neg$ . (3) Applying  $\neg$  on the left and  $\circ$  on the right to (2) we get  $\neg \circ \supseteq \neg \neg \neg \circ$ , whereas applying  $\neg \circ$  on the right to (2) gives  $\neg \circ = \circ \neg \circ \subseteq \neg \neg \neg \circ$ .

(4) Apply  $\setminus$  on the right to (3) and use  $\neg = \circ \setminus$ .

(5) Applying  $\setminus - \setminus$  to (iv) for the last step yields

$$\circ U \cap \neg \neg V = \backslash (- \backslash \circ U \cup \neg V) = \backslash -(\backslash \circ U \cup \neg V)$$
  
=  $\backslash - \backslash (\circ U \cap -V) \subseteq \neg \neg (U \cap V)$ 

(6) The inclusion is clear. So without loss of generality let's assume that U is open. Then  $U = \circ U$  and  $\neg \neg V = \circ \neg \neg V$  such that using (5) twice and then (4) gives

$$-\nabla U \cap \nabla V \subseteq -\nabla U \cap \nabla V = -\nabla U \cap U \cap V = -\nabla U \cap V = -\nabla U \cap U \cap V = -\nabla U \cap U \cap V = -\nabla U \cap U \cap U =$$

As a consequence of the lemma we see using (1) that, if U and V are regular, then so is  $\neg U$  by (3) and  $U \lor V$  by (4) and  $U \land V$  by (6).

Thus  $\neg, \land, \lor$  are indeed well-defined operations on  $\mathcal{R}(X)$ .

We next verify the defining properties of a complete boolean algebra:

(I) Since the partial order  $\leq$  on  $\mathcal{R}(X)$  will be characterized by  $U \leq V \Leftrightarrow U \wedge V = U$ , we must choose  $\leq = \subseteq$ . With this choice it is then obvious that  $\leq$  is a partial order with  $\emptyset$  as a least element and X as a greatest element.

(II) Clearly,  $U \wedge V$  is the infimum of  $\{U, V\}$ . To check the existence of suprema pick  $\mathcal{Y} \subseteq \mathcal{R}(X)$ . As a union of open sets  $Y = \bigcup \mathcal{Y}$  is open. Hence,  $Y \subseteq \neg \neg Y$  by (2) and by (1) and (4)  $\neg \neg Y$  is a regular upper bound of  $\mathcal{Y}$ . Actually,  $\neg \neg Y$  is the supremum of  $\mathcal{Y}$ , since for each regular Z with  $Y \subseteq Z$  we have  $\neg \neg Y \subseteq \neg \neg Z = Z$  by (1).

(III) Both distributive laws for  $\wedge$  and  $\vee$  follow with a straightforward computation from (1) and (6) and the distributive laws for  $\cap$  and  $\cup$ .

(IV) Clearly,  $U \wedge \neg U = U \setminus \overline{U} = \emptyset$ . If U is open, then  $U \cup \neg U = \langle \partial U \rangle$  is a dense set (otherwise there would be an open W with  $W \subseteq \partial U = \overline{U} \setminus U$ , which is absurd). So in particular for regular U we have  $U \vee \neg U = \neg \setminus -(U \cup \neg U) = \neg \setminus X = \neg \emptyset = X$ .

This proves that  $\mathcal{R}(X)$  is a complete boolean algebra.